Addressing

The underlying abstraction is a byte stream addressed by positions from the set:

$$\mathcal{P} = \{ p \mid 0 \leq p \leq \mathtt{size} \} \subset \mathbb{N}_0$$

In slight abuse of notation we define for $s, e \in \mathcal{P}$

$$[s, e] = \begin{cases} \{p \in \mathcal{P} \mid s \le p < e\} & \text{if } s < e \\ s & \text{if } s = e \\ \emptyset & \text{otherwise} \end{cases}$$

Resulting in the set of ranges

$$\mathcal{R} = \{[s, e) \mid s, e \in \mathcal{P}, s < e\}$$

Selections

Selections are non-empty, directed ranges

$$\mathcal{S} = \mathcal{R} \times \{\leftarrow, \downarrow, \rightarrow\}$$

Forming a totally ordered set $\langle S, \leq \rangle$, according to:

$$([s_1, e_1), d_1) \leq ([s_2, e_2), d_2) \iff s_1 < s_2 \lor (s_1 = s_2 \land e_1 \leq e_2)$$

An invalid selection is denoted by (\emptyset, \cdot)

Selections: Membership, Length and Direction

A selection $s = (r, d) \in S$

• contains
$$p \in \mathcal{P} \iff p \in r$$
.

- ▶ has length |s| = |r|
- ▶ is anchored and left extending if $d = \leftarrow$
- ▶ is non-anchored or following if $d = \downarrow$, this implies |s| = 1
- \blacktriangleright is anchored and right extending if $d=\rightarrow$
- can be flipped

$$\mathsf{flip}(s) = egin{cases} (r,
ightarrow) & ext{if } d = \leftarrow \ s & ext{if } d = \downarrow \ (r, \leftarrow) & ext{if } d =
ightarrow \end{cases}$$

Selections: Cursor and Anchor

Each selection has 2 end points, cursor and anchor, which are *over* a character.

$$\mathcal{S} o \mathcal{P}$$

cursor : $([s, e), d) \mapsto \begin{cases} e - 1 & \text{if } d = \to \\ s & \text{otherwise} \end{cases}$
anchor : $([s, e), d) \mapsto \begin{cases} e - 1 & \text{if } d = \leftarrow \\ s & \text{otherwise} \end{cases}$

For singleton selections both cursor and anchor are on the same character.

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Selections: Motions

Normal mode

Non-anchored singleton selections $([p, p+1), \downarrow)$ for cursor at $p \in \mathcal{P}$.

Visual mode

Selections are anchored, only the cursor is adjusted, the anchor remains fixed.

Selections: Cover and Normalization

The cover of $S \subseteq \mathcal{S}$ is

$$\begin{array}{l} \mathsf{cover} \colon 2^{\mathcal{S}} \to 2^{\mathcal{P}} \\ \mathsf{cover}(\mathcal{S}) = \{ p \in \mathcal{P} \mid (r, d) \in \mathcal{S}, p \in r \} \end{array}$$

We define a normalization operator on a set of selections:

normalize:
$$2^{\mathcal{S}} \rightarrow 2^{\mathcal{S}}$$

such that:

1.
$$(\emptyset, \cdot) \notin \text{normalize}(S)$$

2. $\text{cover}(\text{normalize}(S)) = \text{cover}(S)$
3. $\forall s_1 = (r_1, d_1), s_2 = (r_2, d_2) \in S : s_1 \neq s_2 \implies r_1 \cap r_2 = \emptyset$
4. $\forall s_1 = (r_1, d_1) \in S :$
 $(\forall s_2 = (r_2, d_2) \in S \setminus \{s_1\} : r_1 \cap r_2 = \emptyset) \implies s_1 \in \text{normalize}(S)$
A normalized selection set is an interval order.

Selections: Intersect and Union

Operations on ordered sets of selections, $S_1, S_2 \subseteq \mathcal{S}$

intersect

$$S_1 \cap S_2 = \{(s_1 \cap s_2, d_1) \mid (s_1, d_1) \in S_1, (s_2, d_2) \in S_2\}$$

note: not symmetric

union

$$S_1 \cup S_2 = \operatorname{normalize}(S_1 \cup S_2)$$

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Selections: Complement and Minus

- complement
- minus
- These behave as "expected".

TODO: define them formally

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Selections: Pairwise Merge

Pairwise merge with $f:\mathcal{S}\times\mathcal{S}\rightarrow\mathcal{S}$

 $merge(S_1, S_2) = \{f(u_i, v_i) \mid u_1 < \dots < u_n \in S_1, v_1 < \dots < v_m \in S_2\}$

i.e. start with the minimum and repeatedly combine the immediate successor of both sets.

- $\operatorname{union}(([s_1, e_1), d_1), ([s_2, e_2), d_2)) = ([\min(s_1, s_2), \max(e_1, e_2)), d_1)$
- intersect(($(s_1, e_1), d_1$), ($(s_2, e_2), d_2$)) = ($(\max(s_1, s_2), \min(e_1, e_2)), d_1$)

 $\blacktriangleright \operatorname{left}(s_1, s_2) = \min(s_1, s_2)$

$$right(s_1, s_2) = max(s_1, s_2)$$